



Unlocking the Secrets: The Crucial Role of Encoding and Decoding in Cryptography

Payel Mondal¹, Debopriya Dey², Shounak Sinha³ and Krittika Chakraborty⁴

Faculty, Department of Basic Science and Humanities, Narula Institute of Technology, Kolkata, India^{1,2}

Student, Department of Computer Science and Engineering, Narula Institute of Technology, Kolkata, India^{3,4}

Abstract: We live in an information and knowledge economy; learning is not the things you are “done” only in your youth; it must become the digital lifestyle and develop it as a healthy daily habit. When we access bank account online or perform any financial transaction, we need to secure the communication between our computer and the bank's servers. This ensures that our personal and financial information remains confidential and protected from unauthorized access. Moreover, when we send an email, cryptography is used to encrypt the contents of our message so that it can only be read by the intended recipient. This protects the message from being intercepted and read by unauthorized parties. Here cryptology comes into the role to protect sensitive communications and share information privately. This is mainly, the objective of Cryptography, the study of techniques to keep communications private by means of data encryption and its subsequent decryption. Encryption is used to transform an original information (plain form) into some incomprehensible form (coded form).

One commonly used encoding technique is the Advanced Encryption Standard (AES). AES is a symmetric encryption algorithm that uses a single key to both encrypt and decrypt messages. AES is widely used in a variety of applications, including secure file transfers, secure email, and virtual private networks (VPNs). In this paper we are offering an algorithm on encryption and decryption to convert a plain text into a ciphertext and conversely. We used a non-singular matrix as one process of generation of a key to encrypt a message. In this process, the recipient can decode (decrypt) the message by using the inverse of the matrix to get back the original message. But this process is not so effective. For resolving the drawback of this process, we introduced encryption techniques of circular bit rotation algorithm using generating key of random integers of the size number less than 512 bits i.e. at least 154 digits or 77 pairs of unsigned integers. So, the proposed algorithm in this paper encrypts plain text into cipher text that is unrecognizable and which makes the cipher text unidentifiable when compared to plain text.

Keywords: Cryptography, Encryption, Decryption, Plain Text, Cipher Text, ASCII Code, Bit Rotation

I. INTRODUCTION

Encoding and decoding is a crucial part of cryptography. Encoding is the process to transform plain text like letter, words, numbers in a specialized format (ciphertext) using some techniques for effective and safe transformation of data. Decoding is the method of turning back that specialized format again into normal text to receive using some techniques. In this work we will discuss about two kinds of techniques to encode and decode normal message using linear algebra (matrix) and binary operations. We will create "secret" using these techniques and again "unlock that secret".

II. ENCODING AND DECODING TECHNIQUES

Encoding: Plain text has ASCII codes. Arrange all ASCII codes in a form of (2×2) matrix B . Multiply with a generating Matrix or encoding matrix (key) A Plain text again (encoded text).

Decoding: Make the inverse Matrix of encoding matrix. Multiply the matrix A^{-1} and C . Get the decoded matrix B . B is the actual matrix. Transform those ASCII codes into plain text again.

Illustration:

PLAIN TEXT: **Unlock Secrets**

ASCII CODES OF EACH CHARACTER: 85 110 108 111 99 107 32 83 101 99 114 101 99 116
115 46 [ASCII of space=32 and full stop=46]

ENCODING/GENERATING MATRIX(KEY): Let $A = \begin{bmatrix} 1 & 1 & 0 & 1 \end{bmatrix}$, and matrix $B =$
 $\begin{bmatrix} 85 & 108 & 99 & 32 & 101 & 114 & 99 & 115 & 110 & 111 & 107 & 83 & 99 & 101 & 116 & 46 \end{bmatrix}$

Then $C = A * B$

$$C = \begin{bmatrix} 1 & 1 & 0 & 1 \end{bmatrix} * \begin{bmatrix} 85 & 108 & 99 & 32 & 101 & 114 & 99 & 115 & 110 & 111 & 107 & 83 & 99 & 101 & 116 & 46 \end{bmatrix} \\ = \begin{bmatrix} 195 & 119 & 206 & 115 & 200 & 215 & 215 & 161 & 110 & 111 & 107 & 83 & 99 & 101 & 116 & 46 \end{bmatrix}$$



Encoded message = 195 110 119 111 206 107 115 83 200 99 215 101 215 116 161 46

ENCODED PLAIN TEXT [CIPHER TEXT]= $\check{A}nwo\check{i}ksS\check{E}c \times e \times t\check{i}$

DECODING: Take $B = A^{-1} * C$

$$A^{-1} = [1 \ -1 \ 0 \ 1]$$

$$B = [1 \ -1 \ 0 \ 1] * [195 \ 119 \ 206 \ 115 \ 200 \ 215 \ 215 \ 161 \ 110 \ 111 \ 107 \ 83 \ 99 \ 101 \ 116 \ 46]$$

$$= [85 \ 108 \ 99 \ 32 \ 101 \ 114 \ 99 \ 115 \ 110 \ 111 \ 107 \ 83 \ 99 \ 101 \ 116 \ 46]$$

Decoded message: = 85 110 108 111 99 107 32 83 101 99 114 101 99 116 115 46

ENCODED PLAIN TEXT = Unlock Secrets

III. ENCRYPTION AND DECRYPTION TECHNIQUE USING CIRCULAR BIT SHIFT IN BINARY FIELD

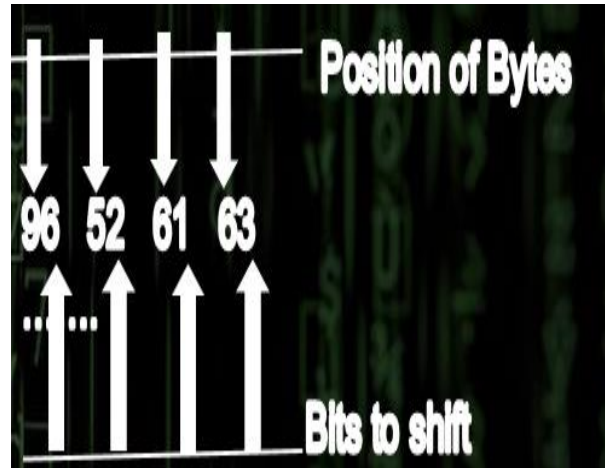
Encryption: Divide plain text into blocks with 10 characters but if there's character number less than 10 in a block use padding instead [1-5].

Generating key:

- Random number generators of size no less than 512 bits i.e. at least 154 digits or 77 pairs of unsigned integers
- Separated into array of pairs and each pair is generated separately
- First digit of each pair is the location of the character and second digit is number of bits to shift as shown in figure

ASCII - Binary Character Table

| Letter | ASCII Code | Binary | Letter | ASCII Code | Binary |
|--------|------------|----------|--------|------------|----------|
| a | 097 | 01100001 | A | 065 | 01000001 |
| b | 098 | 01100010 | B | 066 | 01000010 |
| c | 099 | 01100011 | C | 067 | 01000011 |
| d | 100 | 01100100 | D | 068 | 01000100 |
| e | 101 | 01100101 | E | 069 | 01000101 |
| f | 102 | 01100110 | F | 070 | 01000110 |
| g | 103 | 01100111 | G | 071 | 01000111 |
| h | 104 | 01101000 | H | 072 | 01001000 |
| i | 105 | 01101001 | I | 073 | 01001001 |
| j | 106 | 01101010 | J | 074 | 01001010 |
| k | 107 | 01101011 | K | 075 | 01001011 |
| l | 108 | 01101100 | L | 076 | 01001100 |
| m | 109 | 01101101 | M | 077 | 01001101 |
| n | 110 | 01101110 | N | 078 | 01001110 |
| o | 111 | 01101111 | O | 079 | 01001111 |
| p | 112 | 01110000 | P | 080 | 01010000 |
| q | 113 | 01110001 | Q | 081 | 01010001 |
| r | 114 | 01110010 | R | 082 | 01010010 |
| s | 115 | 01110011 | S | 083 | 01010011 |
| t | 116 | 01110100 | T | 084 | 01010100 |
| u | 117 | 01110101 | U | 085 | 01010101 |
| v | 118 | 01110110 | V | 086 | 01010110 |
| w | 119 | 01110111 | W | 087 | 01010111 |
| x | 120 | 01111000 | X | 088 | 01011000 |
| y | 121 | 01111001 | Y | 089 | 01011001 |
| z | 122 | 01111010 | Z | 090 | 01011010 |



Decryption:

Encoded message in binary \rightarrow circular bit shift operation using the same key but in reverse \rightarrow decoded message \rightarrow decoded plain TEXT BY CONCATENATING

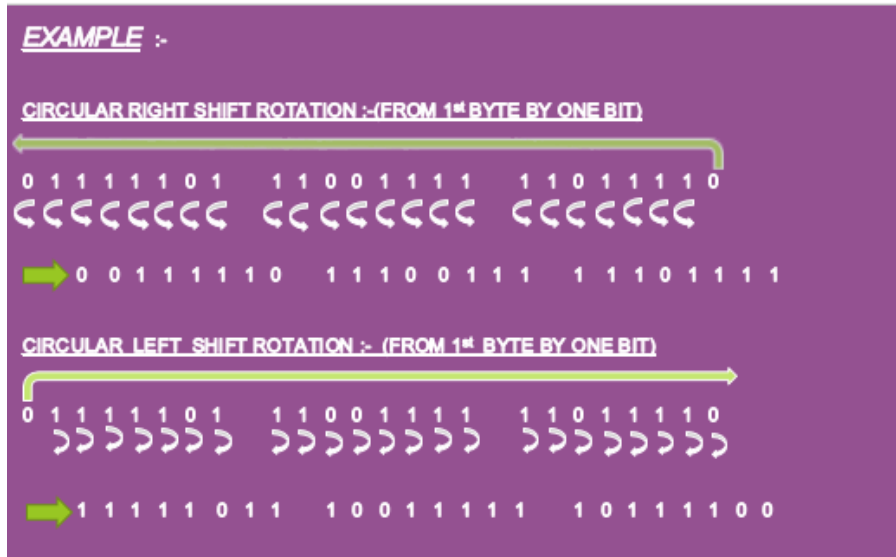
Circular bit shift:

❖ Circular right shift rotation:

In the encryption technique the bits are shifted towards the right and from the end the bits are shifted to the byte from where the bits were shifted.

❖ Circular left shift rotation:

In the decryption technique the bits are shifted towards the left and from the beginning the extra bits are shifted to the byte at the end [6-8]. The position of bytes ranges from 0 to 9 and bits are shifted range from 1 to 8 only, 0 and 9 will invert the bits instead of shifting [9-10].



IV. EXAMPLE

PLAIN TEXT :- Go and Grow
BLOCK 1 :- Go_and_Gro
BLOCK 2 :- w_ _ _ _ _ _ _ _
 For demonstration only 4 pairs of keys are used
KEY :- [24,19,01,00]

i. ENCRYPTION :-

BLOCK 1 :- 01000111 01101111 01011111 01100001 01101110 01100100 01011111 01000111
 01110010 01101111
BLOCK 2 :- 01110111 01011111 01011111 01011111 01011111 01011111 01011111 01011111
 01011111 01011111

ROTATION 1 :-
 Performing a right shift circular rotation on 3rd byte by 4 bits and KEY=24
BLOCK 1 :- 01000111 01101111 11110101 11110110 00010110 11100110 01000101
 11110100 01110111 00100110
BLOCK 2 :- 01101111 01011111 11110101 11110101 11110101 11110101 11110101
 11110101 11110101 11110101

ROTATION 2 :-
 Here the number of bits to shift is 9, so instead of shifting bits will be inverted instead of shifting from 2nd byte.
KEY = 19
BLOCK 1 :- 01000111 10010000 00001010 00001001 11101001 00011001 10111010
 00001011 10001000 11011001
BLOCK 2 :- 01101111 10100000 00001010 00001010 00001010 00001010 00001010
 00001010 00001010 00001010

ROTATION 3 :- Performing a right circular shift rotation on 1st byte by 1 bit
KEY = 01
BLOCK 1 :- 10100011 11001000 00000101 00000100 11110100 10001100 11011101
 00000101 11000100 01101100
BLOCK 2 :- 00111011 11010000 00000101 00000101 00000101 00000101 00000101
 00000101 00000101 00000101

ROTATION 4 :-
 Here the number of bits to shift is 0, so bits will be inverted from 1st to end byte
KEY = 00
BLOCK 1 :- 01011100 00110111 11111010 11111011 00001011 01110011 00100010
 11111010 00111011 10010011



BLOCK 2 :- 11000100 00101111 11111010 11111010 11111010 11111010 11111010
11111010 11111010 11111010

BLOCK 1:- \7úûvTsú”;“

BLOCK 2 :- Ä/úúúúúúúú

BY CONCATENATING

CIPHER TEXT:- \7úû vTsú”;“Ä/úúúúúúúú

ii. DECRYPTION:-

For this purpose, only 4 pairs of keys are used.

KEY:- [24,19,01,00]

BLOCK 1 :- 01011100 00110111 11111010 11111011 00001011 01110011
00100010 11111010 00111011 10010011

BLOCK 2 :- 11000100 00101111 11111010 11111010 11111010 11111010 11111010
11111010 11111010 11111010

ROTATION 1 :-

Here the number of bits to shift are 0 ,so bits will be inverted from 1st byte to end end byte

KEY :- 00

BLOCK 1:- 10100011 11001000 00000101 00000100 11110100 10001100 11011101
00000101 11000100 01101100

BLOCK 2:- 00111011 11010000 00000101 00000101 00000101 00000101
00000101 00000101 00000101 00000101

ROTATION 2 :-

Performing a left circular shift rotation on 1st byte by 1 bit

KEY:- 01

BLOCK 1 :- 01000111 10010000 00001010 00001001 11101001 00011001 10111010
00001011 10001000 11011001

BLOCK 2:- 01110111 10100000 00001010 00001010 00001010 00001010 00001010
00001010 00001010 00001010

ROTATION 3 :-

Here the number of bits to shift is 9, so, instead of shifting bits will be inverted instead of shifting from 2nd byte.

KEY = 19

BLOCK 1 :- 01000111 01101111 11110101 11110110 00010110 11100110 01000101
11110100 01110111 00100110

BLOCK 2 :- 01110111 01011111 11110101 11110101 11110101 11110101 11110101
11110101 11110101 11110101

ROTATION 4 :-

Performing a right shift circular rotation on 3rd byte by 4 bits

KEY: 24

BLOCK 1:- 01000111 01101111 01011111 01100001 01101110 01100100 01011111
01000111 01110010 01101111

BLOCK 2:- 01110111 01011111 01011111 01011111 01011111 01011111 01011111
01011111 01011111 01011111

DECODED MESSAGE:-

BLOCK 1 :- GoandGro

BLOCK 2 :- w_____

BY CONCATENATING:- *-Go and Grow*

V. CONCLUSION

Cryptography helps us to secure our data privacy and it’s not unknown how data is important in our daily life and for organizations. These two techniques are many of those to ensure data security and mathematics is playing a huge role in it. But these algorithms have drawbacks also. So, there are further scopes to enhance the algorithms by increasing efficiency and reducing time complexity. Again, there can be error during message transmission for a lot of reason. We have to enhance our algorithm always to detect and correct errors more easily and efficiently.



REFERENCES

- [1]. Bhardwaj A. and Som S. (2016). Study of different cryptographic technique and challenges in future. 2016 International Conference on Innovation and Challenges in Cyber Security (ICICCS-INBUSH), Noida, 208-212.
- [2]. Kunwar V., Agarwal N., Rana A. and Pandey J. P. (2018). Load balancing in cloud—a systematic review. *Advances in Intelligent Systems and Computing*, 654, 583-593.
- [3]. Chauhan B. D., Rana A. and Sharma N. (2017). Testing sufficiency test (TST) - Evolving a new model for estimating software test cases. *International Journal of Applied Engineering Research*, 12-21.
- [4]. Dubey G., Rana A. and Ranjan J. (2017). Fine-grained opinion mining of product review using sentiment and semantic orientation. *International Journal of Business Information Systems*, 25 (1), 1-17.
- [5]. Ghosh S., Rana A. and Kansal V. (2017). Predicting defect of software system. *Advances in Intelligent Systems and Computing*, 516, 55- 67.
- [6]. Nath S. and Som S. (2017, January). Security and Privacy Challenges: Internet of Things. *Indian Journal of Science and Technology*, 10(3), DOI: 10.17485/ijst/2017/v10i3/110642.
- [7]. Goldwasser S. and Micali S. (1984, April). Probabilistic encryption. *J. Comput. Syst. Sci.* 28(2), 270-299.
- [8]. Rivest R. L., Shamir A. and Adleman L. M. (1983): A method for obtaining digital signatures and public-key cryptosystems (reprint). *Commun. ACM* 26(1).
- [9]. El Gamal T. (1985). A public key cryptosystem and a signature scheme based on discrete logarithms. In: Blakely, G.R., Chaum, D. (eds.) *CRYPTO 1984*. LNCS, 196, 10-18.
- [10]. Uludag U., Pankanti S., Prabhakar S. and Jain A. K. (2004, June). Biometric cryptosystems: Issues and challenges. *Proc. IEEE (Special Issue Multimedia Security for Digital Rights Management)*, 92 (6), 948 -960.